

Note

A relative of the Thue–Morse sequence

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Abstract

We study a sequence, c , which encodes the lengths of blocks in the Thue–Morse sequence. In particular, we show that the generating function for c is a simple product.

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Consider the sequence

$$c: c_0, c_1, c_2, c_3, \dots = 1, 3, 4, 5, 7, 9, 11, 12, 13, \dots$$

defined to be the lexicographically least sequence of positive integers satisfying $n \in c$ implies $2n \notin c$. In fact, the lexicographic minimality of c makes it possible to replace the previous ‘implies’ with ‘if and only if.’ Equivalently, c is defined inductively by $c_0 = 1$ and

$$c_{k+1} = \begin{cases} c_k + 1 & \text{if } (c_k + 1)/2 \notin c, \\ c_k + 2 & \text{otherwise,} \end{cases} \quad (1)$$

for $k \geq 0$. This sequence was the focus of a problem of Kimberling in [7]. (In fact, he looked at the sequence $4c_0, 4c_1, 4c_2, \dots$.) The solution was given by Bloom [4]. Our Corollary 7 answers essentially the same question. Related results have recently been announced by Tamura [9].

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At the 4^e Colloque Séries Formelles et Combinatoire Algébrique (Montréal, June 1992) Plouffe and Zimmermann [8] posed the following problem. Show that the generating function for c is

$$\sum_{k \geq 0} c_k x^k = \frac{1}{1-x} \prod_{j \geq 1} \frac{1-x^{2e_j}}{1-x^{e_j}} = \frac{1}{1-x} \prod_{j \geq 1} (1+x^{e_j}) \quad (2)$$

the sequence of exponents being

$$e: e_1, e_2, e_3, e_4, \dots = 1, 1, 3, 5, 11, 21, 43, \dots,$$

where $e_1 = 1$ and

$$e_{j+1} = \begin{cases} 2e_j + 1 & \text{if } j \text{ is even,} \\ 2e_j - 1 & \text{if } j \text{ is odd,} \end{cases} \quad (3)$$

for $j \geq 1$. They found this conjecture by using a method that goes back to Euler. First they assumed that the generating function was of the form

$$\prod_{j \geq 0} \frac{1-x^{a_j}}{1-x^{b_j}}$$

for a certain pair of sequences a_j, b_j . Then they took the logarithm to convert the product into a sum. Finally they used Möbius inversion to determine the candidate sequences. Details of this procedure can be found in the text of Andrews [2, Theorem 10.3].

The purpose of this note is to prove (2). Before doing this, however, we will show that c has a number of other interesting properties. Chief among these is the fact that c is closely related to the famous Thue–Morse sequence, t . See the survey article of Berstel [3] for more information about t .

First we need to have a characterization of the integers in the sequence c .

Proposition 1. *If n is any positive integer then $n \in c$ if and only if $n = 2^{2^i}(2j+1)$ for some nonnegative integers i and j .*

Proof. Every positive integer n can be uniquely written in the form $n = 2^k(2j+1)$ where $k, j \geq 0$. We will proceed by induction on k .

If $k=0$, then n is odd. But then $n/2$ is not an integer, and so n is in the sequence by definition (1).

Now assume that $k \geq 1$ and that the proposition holds for all powers less than k of 2. If $k=2i$ is even, then by induction we have $2^{2^{i-1}}(2j+1) \notin c$. So $n = 2^{2^i}(2j+1) \in c$ by (1). On the other hand, if $k=2i+1$ is odd, then induction implies $2^{2^i}(2j+1) \in c$. Thus $n = 2^{2^{i+1}}(2j+1) \notin c$ as desired. \square

Let χ be the characteristic function of c , i.e.,

$$\chi(n) = \begin{cases} 1 & \text{if } n \in c, \\ 0 & \text{otherwise.} \end{cases}$$

Restating the previous proposition in terms of χ yields the next result.

Lemma 2. *The function χ is uniquely determined by the equations*

$$\chi(2n + 1) = 1$$

$$\chi(4n + 2) = 0$$

$$\chi(4n) = \chi(n).$$

Another way of obtaining the sequence $\chi(n)$ for $n \geq 1$ is as follows. Starting from the sequence

$$101 \bullet 101 \bullet 101 \bullet 101 \bullet \dots$$

defined on the alphabet $\{0, 1, \bullet\}$, fill in the successive holes with the successive terms of the sequence itself, obtaining:

$$101110101011101 \bullet \dots$$

Iterating this process infinitely many times (by inserting the initial sequence into the holes at each step), one gets a ‘Toeplitz transform’ which is nothing but our sequence χ . The proof of this fact is easily obtained using Lemma 2. See the article of Allouche and Bacher [1] for more information about Toeplitz transformations.

The connection with the Thue–Morse sequence can now be obtained. This sequence is

$$t: t_0, t_1, t_2, t_3, \dots = 0, 1, 1, 0, 1, 0, 0, 1, \dots$$

defined by the conditions

$$t_0 = 0,$$

$$t_{2n+1} \equiv t_n + 1 \pmod{2},$$

$$t_{2n} = t_n.$$

We will need a lemma relating t and χ . All congruences in this and any future results will be mod 2.

Lemma 3. *For every positive integer, n , we have*

$$\chi(n) \equiv t_n + t_{n-1}.$$

Proof. This is a three case induction based on Lemma 2 and the definitions of χ and t . We will only do one of the cases as the others are similar.

$$\begin{aligned} t_{4n} + t_{4n-1} &\equiv t_{2n} + t_{2n-1} + 1 \\ &\equiv t_n + t_{n-1} + 2 \\ &\equiv \chi(n) \\ &= \chi(4n). \quad \square \end{aligned}$$

Define d_k to be the first difference sequence of c_k , i.e., $d_k = c_k - c_{k-1}$, for $k \geq 0$ ($c_{-1} = 0$). So d is the sequence

$$d_0, d_1, d_2, d_3, d_4, \dots = 1, 2, 1, 1, 2, 2, 2, 1, 1, 2, 1, \dots$$

Note that from the definition of c in (1), the value of d_k is either 1 or 2. Write the Thue–Morse sequence in term of its blocks

$$t = 011010011 \dots = 0^{d_0} 1^{d_1} 0^{d_2} 1^{d_3} \dots$$

defining a sequence d'_k . It is this sequence that is related to our original one via the difference operator.

Theorem 4. For all $k \geq 0$ we have $d_k = d'_k$.

Proof. Since both sequences consist of 1's and 2's, we need only verify that the 1's appear in the same places in both. It will be convenient to let $c'_k = \sum_{i \leq k} d'_i$. We now proceed by induction on k , assuming that $d_i = d'_i$ for $i \leq k$. Then, from the definitions,

$$d_{k+1} = 1 \Leftrightarrow \chi(c_k + 1) = 1. \quad (4)$$

But by the induction hypothesis, $c_k = \sum_{i \leq k} d_i = \sum_{i \leq k} d'_i = c'_k$. So, from Eq. (4),

$$\begin{aligned} d_{k+1} = 1 &\Leftrightarrow \chi(c'_k + 1) = 1 \\ &\Leftrightarrow t_{c'_k+1} + t_{c'_k} \equiv 1 \quad (\text{Lemma 3}) \\ &\Leftrightarrow t_{c'_k+1} \neq t_{c'_k} \\ &\Leftrightarrow d'_{k+1} = 1 \quad (\text{definitions}). \quad \square \end{aligned}$$

Brek [5] used the sequence d in calculating the number of factors of t of given length. The paper of de Luca and Varricchio [6] attacks the same problem in a different way.

Now if $n \in c$ then we will consider its rank, $r(n)$, which is the function satisfying $c_{r(n)} = n$. Note that $r(n)$ is not defined for all positive integers n . In order to obtain a formula for $r(n)$, we will need a definition. Let the base 2 expansion of n be

$$n = \sum_{i \geq 0} \varepsilon_i 2^i$$

with the $\varepsilon_i \in \{0, 1\}$ for all i . Define a function s by

$$s(n) = \sum_{i \geq 0} (-1)^i \varepsilon_i.$$

In other words, $s(n)$ is the alternating sum of the binary digits of n .

Theorem 5. *If $n \in c$ then*

$$r(n) = (2n + s(n))/3 - 1. \tag{5}$$

Proof. The proof will be by induction. From Proposition 1, $n \in c$ if and only if n is odd or $n = 2^{2i}(2j + 1)$ where $i > 0$ and $j \geq 0$. To facilitate the induction, it will be convenient to split the odd numbers into two groups depending upon whether the highest power of 2 dividing $n + 1$ is even or odd. So there will be three cases

- (1) $n = 2^{2i}(2j + 1)$,
- (2) $n = 2^{2i}(2j + 1) - 1$,
- (3) $n = 2^{2i-1}(2j + 1) - 1$,

where $i > 0$ and $j \geq 0$. The arguments are similar, so we will only do the first case.

So suppose n is even (remember that $i > 0$). Thus $n + 1$ is odd and, by Proposition 1, we have $n + 1 \in c$. Since both n and $n + 1$ are in c , the left-hand side of Eq. (5) satisfies $r(n + 1) = r(n) + 1$. So, by induction, it suffices to show that $r'(n + 1) = r'(n) + 1$ where $r'(n)$ is the right-hand side of this equation. Moreover, n is a multiple of 4, hence $s(n + 1) = s(n) + 1$ (write down their binary expansions). Thus

$$\begin{aligned} r'(n + 1) &= (2n + 2 + s(n + 1))/3 - 1 \\ &= (2n + 2 + s(n) + 1)/3 - 1 \\ &= (2n + s(n))/3 \\ &= r'(n) + 1. \quad \square \end{aligned}$$

As a straightforward corollaries we have the next two results.

Corollary 6. *If $n \in c$ then*

$$r(n) = 2n/3 + O(\log n)$$

and $r(n)$ takes the value $2n/3$ infinitely often.

Corollary 7. *For any nonnegative integer k*

$$c_k = 3k/2 + O(\log k)$$

and $c_k = 3k/2$ infinitely often.

We shall now prove the identity (2). First we note a property of the exponents e_j which is a simple consequence of their definition (3).

Lemma 8. For $k \geq 2$, let $f_k = \sum_{2 \leq j \leq k} e_j$. Then

$$f_k = \begin{cases} e_{k+1} - 2 & \text{if } k \text{ is even,} \\ e_{k+1} - 1 & \text{if } k \text{ is odd.} \end{cases}$$

Finally, we come to the proof. We restate the generating function here for easy reference.

Theorem 9. The generating function for c is

$$\sum_{k \geq 0} c_k x^k = \frac{1}{1-x} \prod_{j \geq 1} (1 + x^{e_j}).$$

Proof. It suffices to show that if $k \geq 2$ then

$$g_k(x) = \frac{1}{1-x} (1+x^1)(1+x^1)(1+x^3) \cdots (1+x^{e_k})$$

is the generating function for the sequence

$$1, 3, 4, 5, 7, \dots, c_{f_k}, 2^k, 2^k, 2^k, \dots$$

with $c_{f_k} = 2^k - 1$. The proof is an induction, breaking up into two parts depending on the parity of k . We will do the case where k is odd. (Even k is similar.) Now, by Lemma 8, $g_k(x)(1+x^{e_{k+1}})$ is the generating function for the sequence

$$1, 3, \dots, c_{f_k}, 2^k + 1, 2^k + 3, \dots, 2^k + c_{f_k}, 2^{k+1}, 2^{k+1}, \dots$$

Using Proposition 1 and the fact that k is odd, we see that $2^k + 1 = c_{f_{k+1}}$ and $2^k + c_{f_k} = 2^{k+1} - 1 = c_{f_{k+1}}$. So we want to show that

$$c_{f_{k+1}}, c_{f_{k+2}}, \dots, c_{f_{k+1}} = 2^k + c_0, 2^k + c_1, \dots, 2^k + c_{f_k}.$$

But if $n < 2^k$, then the highest power of 2 dividing n is equal to the highest power dividing $2^k + n$. Thus, by Proposition 1 again, $n \in c$ if and only if $2^k + n \in c$. This gives us the desired equality of the two sequences. \square

One possible generalization of c is the sequence $c^{(\alpha)}$ defined by $n \in c^{(\alpha)}$ if and only if $\alpha n \notin c^{(\alpha)}$. Thus c is the special case $\alpha = 2$.

The following observation is a direct consequence of our definitions.

Proposition 10. *If $\chi^{(\alpha)}(n)$ is the characteristic function of $e^{(\alpha)}$, then the sequence $(\chi^{(\alpha)}(n))$ is the unique fixed point of the morphism*

$$\begin{aligned} 1 &\rightarrow 1^{x-1}0, \\ 0 &\rightarrow 1^{x-1}1 \end{aligned}$$

which begins with 1.

One can also see that $e^{(\alpha)}$ satisfies analogs of many of our previous theorems. For example, if one defines $e_1^{(\alpha)} = 1$ and

$$e_{j+1}^{(\alpha)} = \begin{cases} \alpha e_j^{(\alpha)} + 1 & \text{if } j \text{ is even,} \\ \alpha e_j^{(\alpha)} - 1 & \text{if } j \text{ is odd,} \end{cases}$$

for $j \geq 1$, then the following result is a generalization of Theorem 9 and has an analogous proof.

Theorem 11. *The generating function for $e^{(\alpha)}$ is*

$$\frac{1}{1-x} \prod_{j \geq 1} \frac{1 - x^{\alpha e_j^{(\alpha)}}}{1 - x^{e_j^{(\alpha)}}}.$$

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